



Hayroll: HARVEST Annotator for Yielding Regions Of Lexical Logic A Modular Wrapper for Translating C Macros and Conditional Compilation to Rust



Haoran Peng, Baris Kasikci, Gilbert Louis Bernstein, and Michael D. Ernst
University of Washington

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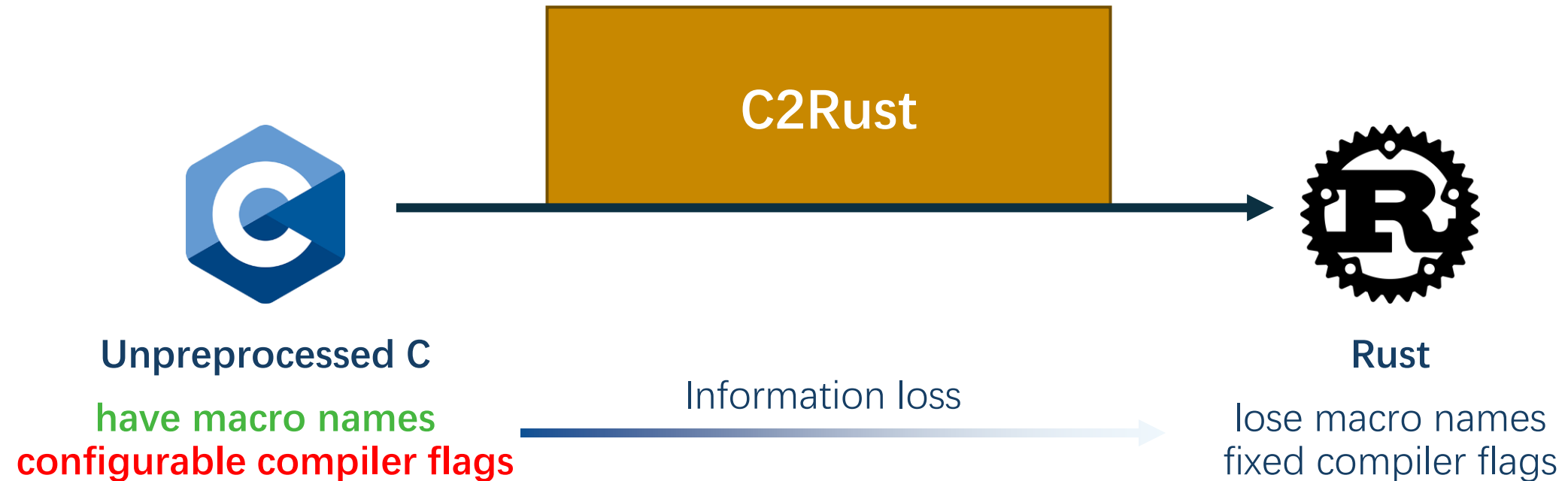


Why Translate C to Rust

- Skipped. Thanks to Cpp2Rust and &inator.
- The C programming language underlies most of the world's software vulnerabilities, mostly due to **memory unsafety**
- About 70% of critical vulnerabilities in Android and Chrome are due to **memory unsafety**
- Rust is designed to be a **memory-safe** language
- A White House report mentions **Rust as a memory-safe** alternative to C/C++
 - Back to the Building Blocks: A Path toward Secure and Measurable Software, Feb 2024, The White House

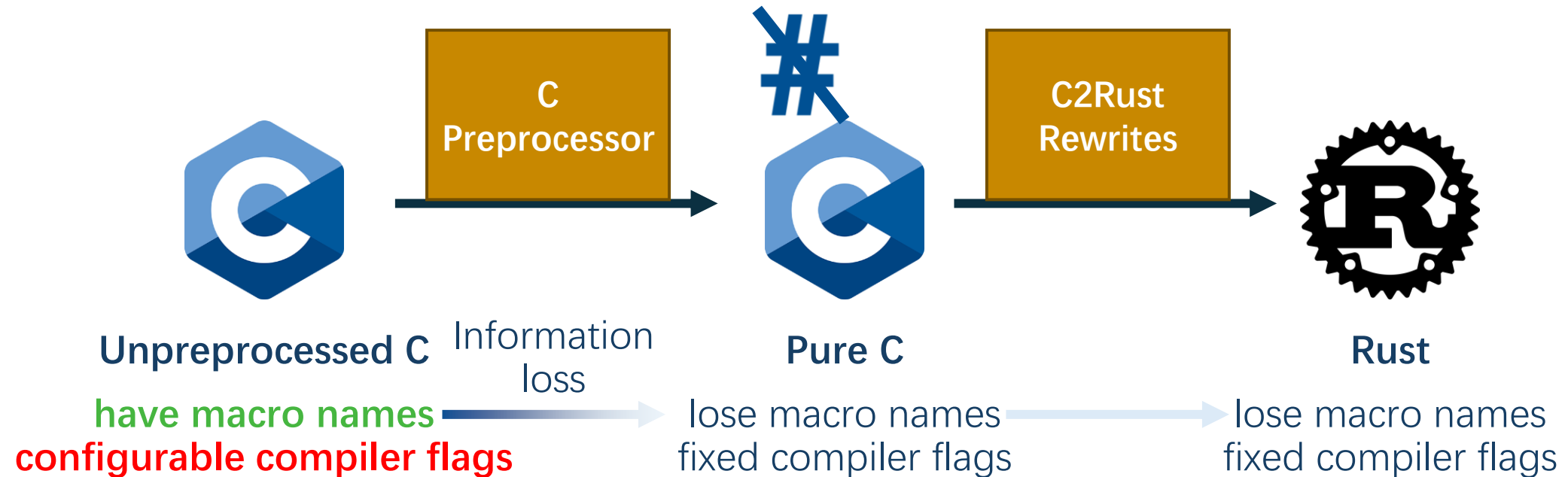


Existing Transpilers Lose **Macros** and **Configurability**





Existing Transpilers Run the C Preprocessor in Advance



Macros and Conditional Compilation is an Integral Part of C



- Hayroll translates **C macros** into **Rust functions or macros**
- Hayroll translates **C conditional compilation** into **Rust cfg attr**
- Hayroll is the first and, as of now, the only system that handles this problem



The C Preprocessor Expands Macros and Concretizes Conditional Compilation



```
float sinh(float x)
{
#ifdef __LIBMCS_FPU_DAZ
    x *= __volatile_onef;
#endif
    float t, w, h;
    int32_t ix, jx;
    GET_FLOAT_WORD(jx, x);
    ix = jx & 0x7fffffff;
    /* x is INF or NaN */
    if (!FLT_UNWORD_IS_FINITE(ix)) {
        return x + x;
    }

    h = 0.5f;
```



```
float sinh(float x)
{
    float t, w, h;
    int32_t ix, jx;
    do {
        ieee_float_shape_type gf_u;
        gf_u.value = (x);
        (jx) = gf_u.word;
    } while (0 == 1);
    ix = jx & 0x7fffffff;
    /* x is INF or NaN */
    if (!((ix) < 0x7f800000L)) {
        return x + x;
    }

    h = 0.5f;
```

C2Rust Starts from Preprocessed C Code



Cond Stmt Expr



```
float sinh(float x)
{
#ifdef __LIBMCS_FPU_DAZ
    x *= __volatile_onef;
#endif
    float t, w, h;
    int32_t ix, jx;
    GET_FLOAT_WORD(jx, x);
    ix = jx & 0x7fffffff;
    /* x is INF or NaN */
    if (!FLT_UNWORD_IS_FINITE(ix)) {
        return x + x;
    }

    h = 0.5f;
```

Challenge 1:
Source correspondence

Challenge 2:
Configuration explosion



-DLIBMCS_FPU_DAZ

```
pub unsafe extern "C" fn sinh(mut x: libc::c_float) -> libc::c_float {
    let mut x = x;
    x *= __volatile_onef;
    let mut t: libc::c_float = 0.;
    let mut w: libc::c_float = 0.;
    let mut h: libc::c_float = 0.;
    let mut ix: int32_t = 0;
    let mut jx: int32_t = 0;
    loop {
        let mut gf_u = ieee_float_shape_type { value: 0. };
        gf_u.value = x;
        jx = gf_u.word as int32_t;
        if !(0 as libc::c_int == 1 as libc::c_int) {
            break;
        }
    }
    ix = jx & 0x7fffffff as libc::c_int;
    if !((ix as libc::c_long) < 0x7f800000 as libc::c_long) {
        return x + x;
    }
}
```



Hayroll Preserves Macros and Conditional Compilation into Rust

Cond Stmt Expr



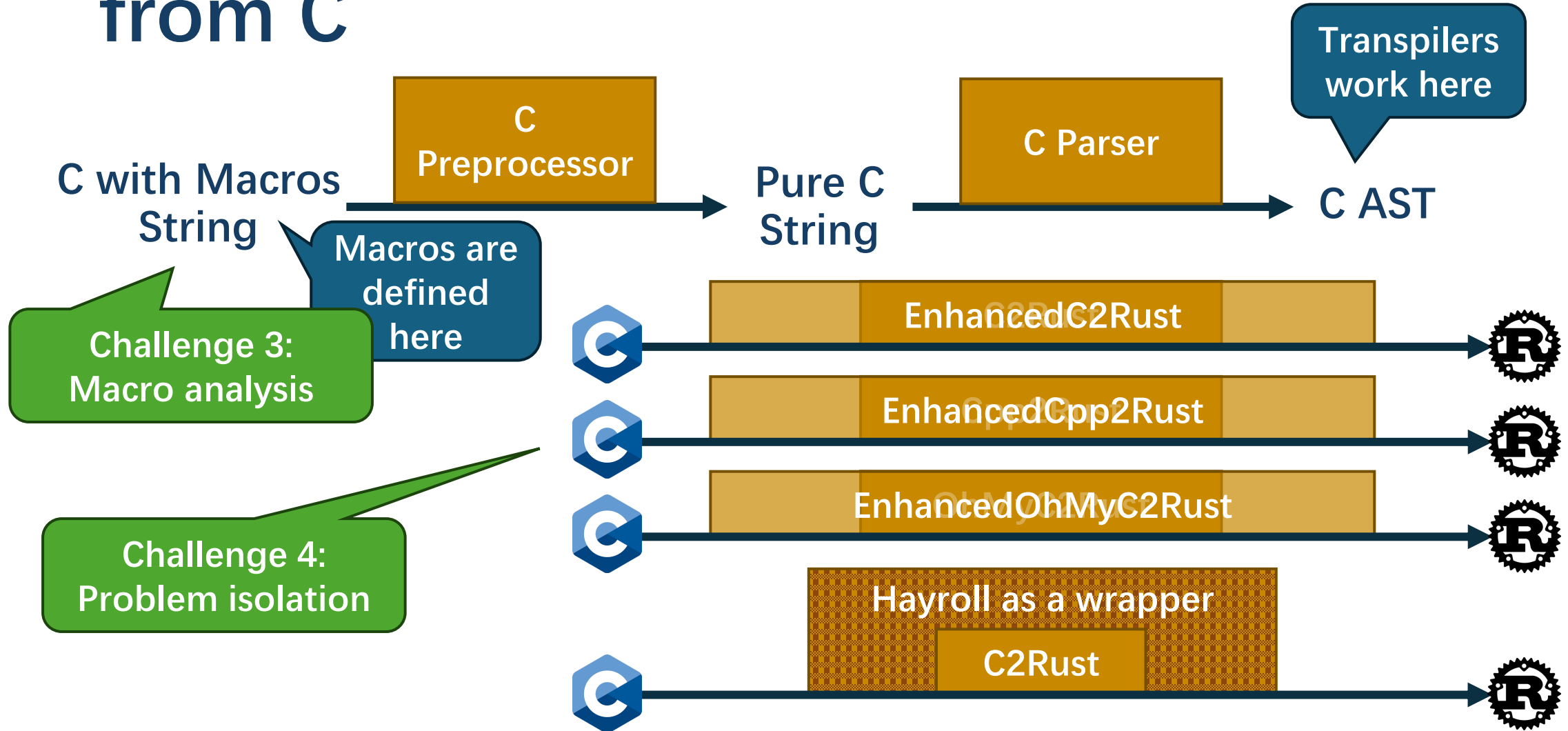
```
float sinh(float x)
{
#ifdef __LIBMCS_FPU_DAZ
    x *= __volatile_onef;
#endif
    float t, w, h;
    int32_t ix, jx;
    GET_FLOAT_WORD(jx, x);
    ix = jx & 0x7fffffff;
    /* x is INF or NaN */
    if (!FLT_UWORD_IS_FINITE(ix)) {
        return x + x;
    }

    h = 0.5f;
```



```
pub unsafe extern "C" fn sinh(mut x: libc::c_float) -> libc::c_float {
    #[cfg(feature = "LIBMCS_FPU_DAZ")]
    (x *= __volatile_onef);
    let mut t: libc::c_float = 0.;
    let mut w: libc::c_float = 0.;
    let mut h: libc::c_float = 0.;
    let mut ix: int32_t = 0;
    let mut jx: int32_t = 0;
    GET_FLOAT_WORD(&mut jx, &mut x);
    ix = jx & 0x7fffffff as libc::c_int;
    if FLT_UWORD_IS_FINITE(&mut ix as *mut int32_t) == 0 {
        return x + x;
    }
    h = 0.5f32;
```


C Preprocessor is a Different Language from C



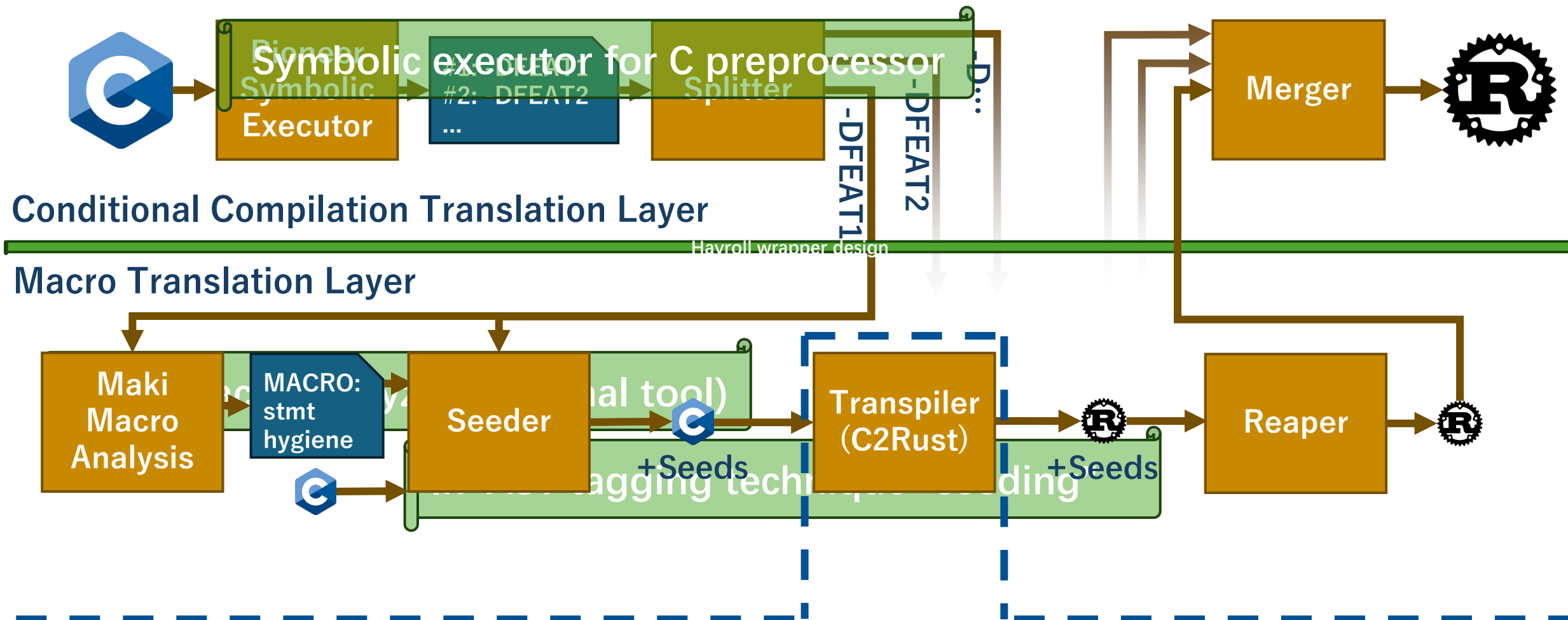


Challenges and Hayroll's Solutions

Challenge	Solution
Problem isolation	Hayroll wrapper design
Macro analysis	Maki macro analyzer (external tool)
Configuration explosion	Symbolic executor for C preprocessor
Source correspondence	In-AST tagging technique “seeding”



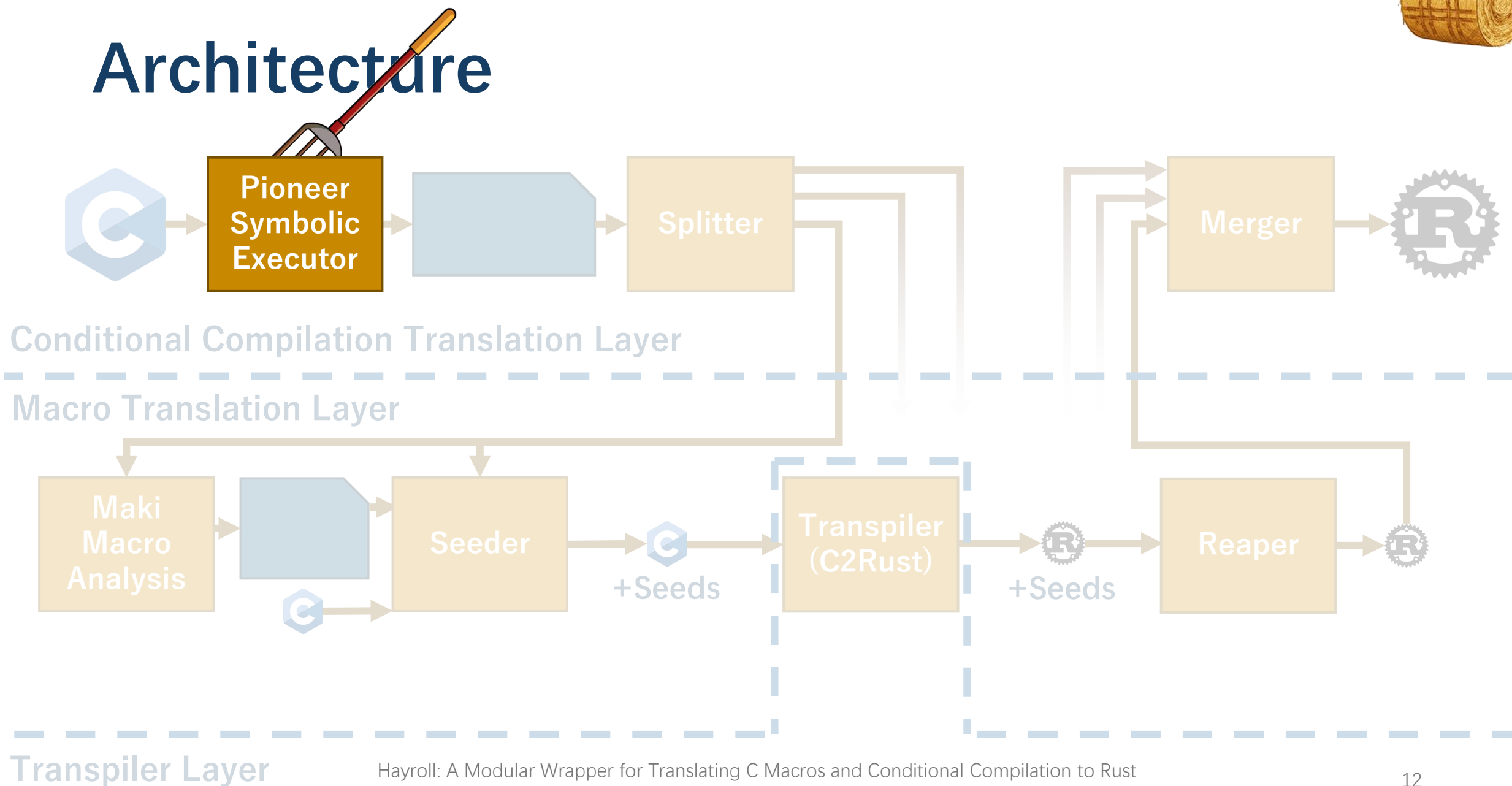
Architecture



Transpiler Layer



Architecture





Symbolic Executor: C Conditional Compilation Requires Global Reasoning

To enable this conditional compilation, shall we just define `__LIBMCS_FPU_DAZ`?



```
float sinh(float x)
{
#ifdef __LIBMCS_FPU_DAZ
    x *= __volatile_onef;
#endif
}
```

gcc -D__LIBMCS_FPU_DAZ ...
Error: undefined reference to
__volatile_onef



Symbolic Executor: C Conditional Compilation Requires Global Reasoning

The user should define `LIBMCS_FPU_DAZ` to trigger **underscored** `__LIBMCS_FPU_DAZ`



```
// from #include "internal_config.h"
#ifdef LIBMCS_FPU_DAZ
    #define __LIBMCS_FPU_DAZ
    static volatile double __volatile_one = 1.0;
    static volatile float __volatile_onef = 1.0f;
#endif

float sinh(double x)
{
    #ifdef __LIBMCS_FPU_DAZ
        x *= __volatile_onef;
    #endif
}
```

Symbolic Executor output:

- internal_config.h:129:5~135:1
requires

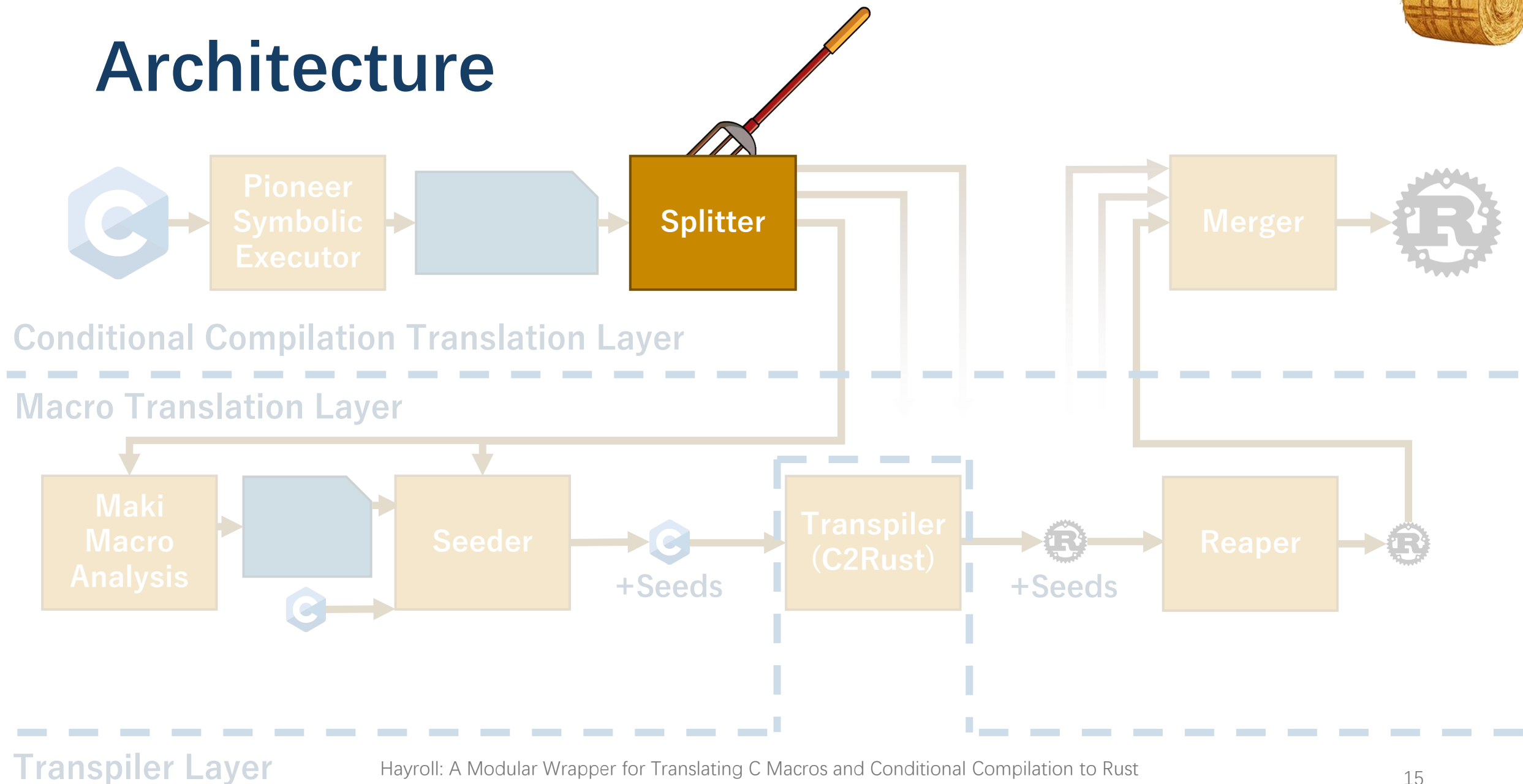
defLIBMCS_FPU_DAZ

- sinh.c:13:5~13:26
requires

defLIBMCS_FPU_DAZ



Architecture





Splitter: Splitting a Configurable Program

into many configuration-specific programs that **together cover every line of code**



```
#ifdef FEAT1
#   ifdef FEAT2
#   endif
#endif
#ifdef FEAT2
#endif
#ifdef FEAT3
#endif
#ifndef FEAT3
#endif
```

All possible combinations

- (no -D flags)
- -DFEAT1
- -DFEAT2
- -DFEAT3
- -DFEAT1 -DFEAT2
- -DFEAT1 -DFEAT3
- -DFEAT2 -DFEAT3
- -DFEAT1 -DFEAT2 -DFEAT3



Splitter: Splitting a Configurable Program

into many configuration-specific programs that **together cover every line of code**



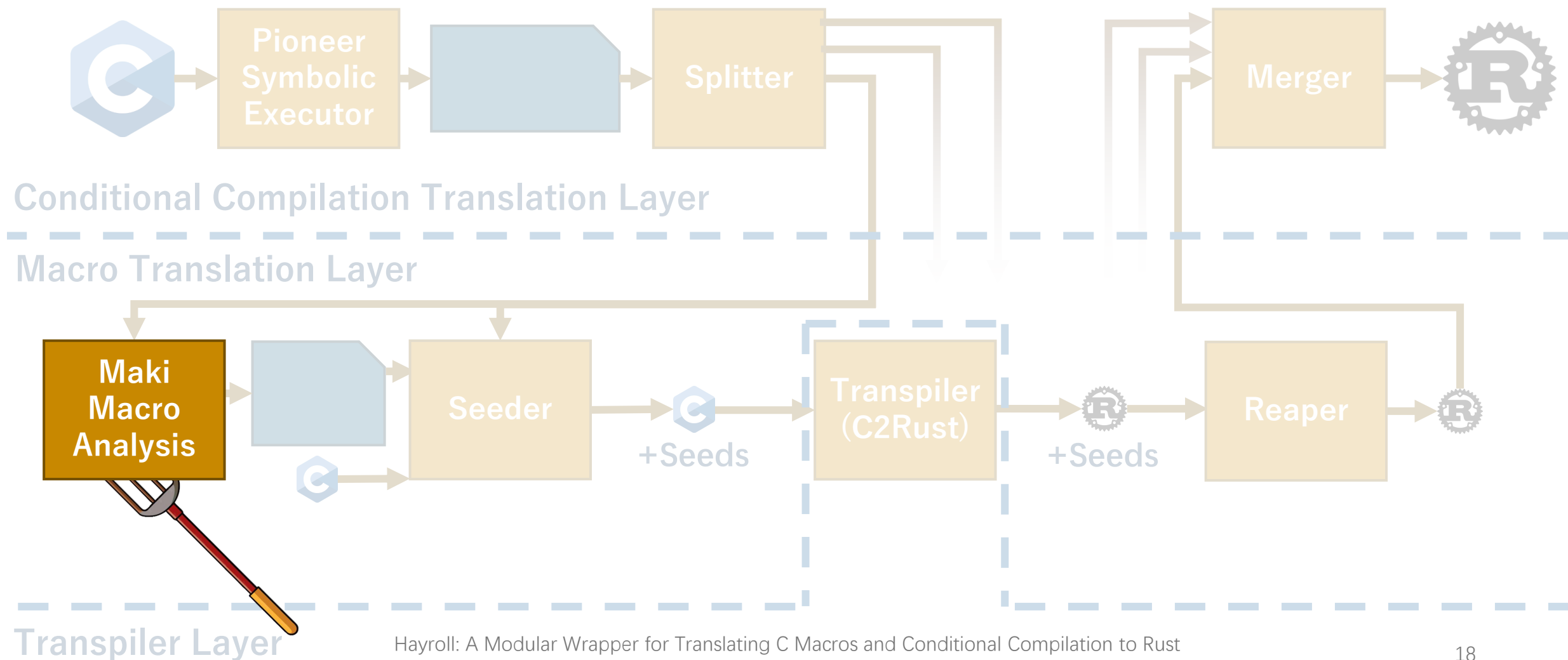
```
#ifdef FEAT1
#   ifdef FEAT2
#   endif
#endif
#ifdef FEAT2
#endif
#ifdef FEAT3
#endif
#ifndef FEAT3
#endif
```

By greedy algorithm
(deepest first)

- **-DFEAT1 -DFEAT2**
- **-DFEAT3**
- **(no -D flags)**



Architecture





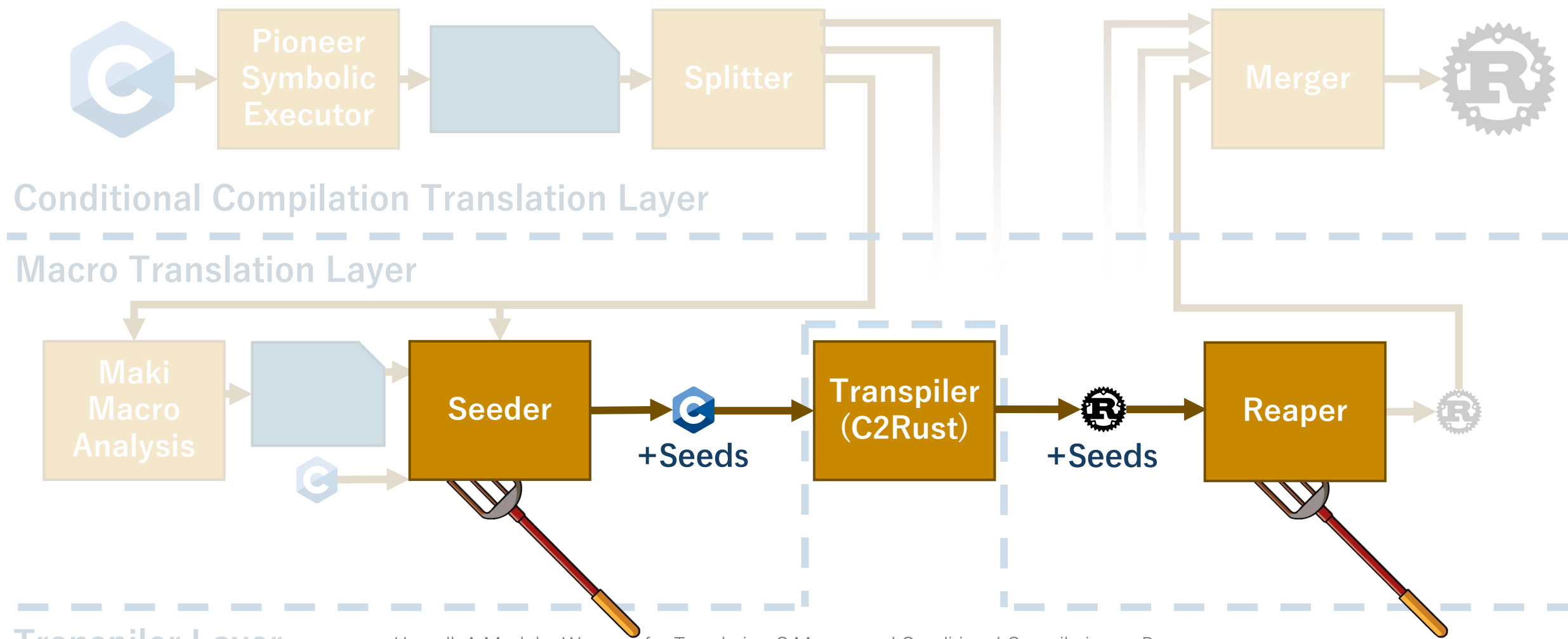
Macro Translatability Properties

- **Leave expanded**
 - Non-syntactical (#define semicolon ;)
 - Non-hygienic (captures local vars)
 - Uses pasting (id##_suffix -> id_suffix)
 - Uses stringification (#name -> "name")
 - Has function pointer
 - Uses integer constant expression (int a[**N**];)
- **Translate into Rust macro**
 - None of the above properties
 - Has control flow jumps (return, continue, break, goto)
 - Has non-expr arguments (MACRO_FOR(3, {printf(...)}))
 - Argument has side effect (M(i++))
- **Translate into Rust function**
 - None of the above properties

Maki (Hayroll-enhanced version) captures these properties by implementing hooks in Clang's C preprocessor



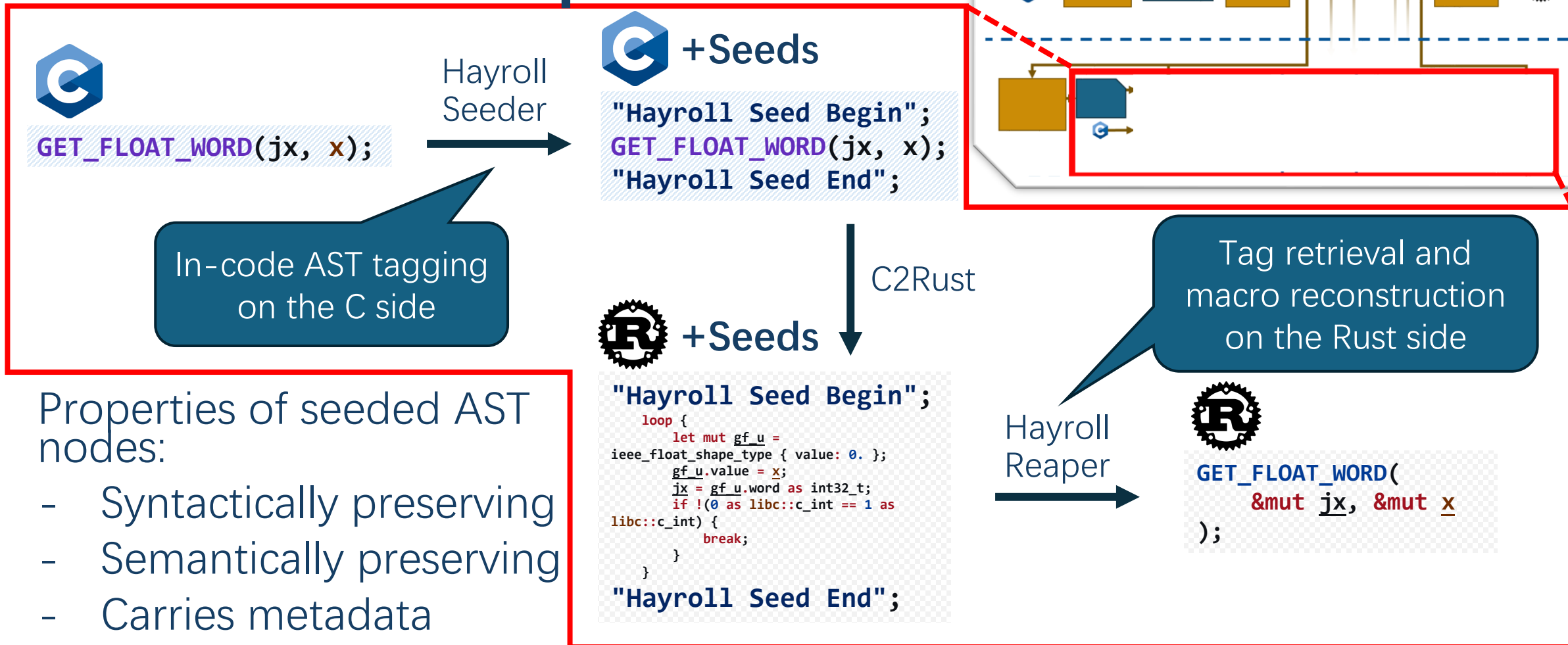
Architecture



Transpiler Layer

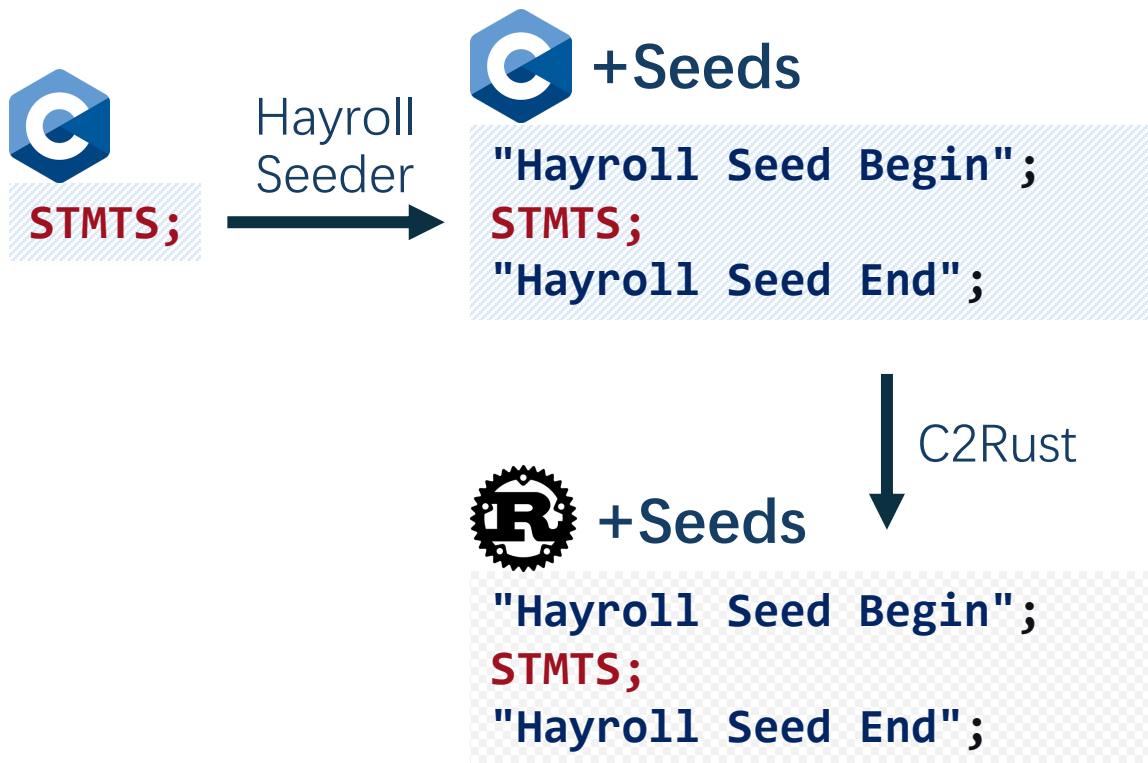


Seeder+Reaper for Source Correspondence



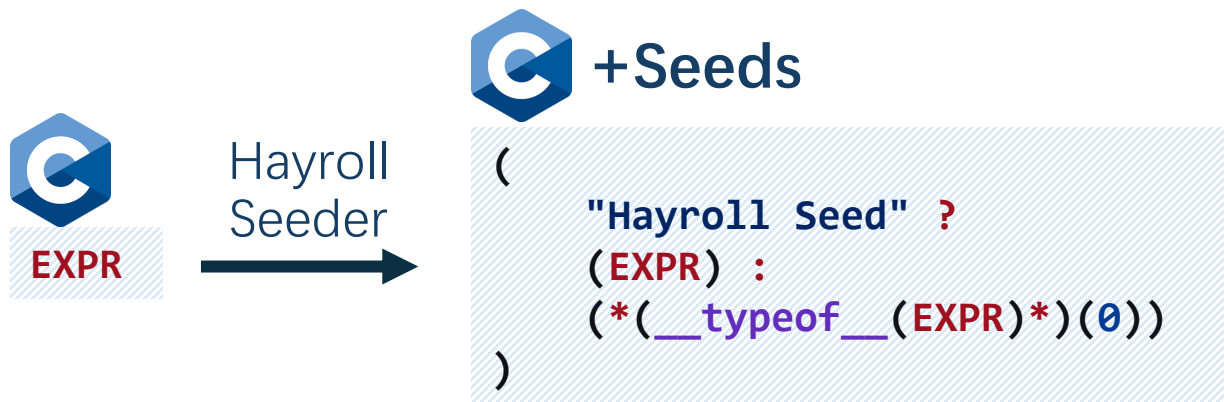


Seeding Statement(s)



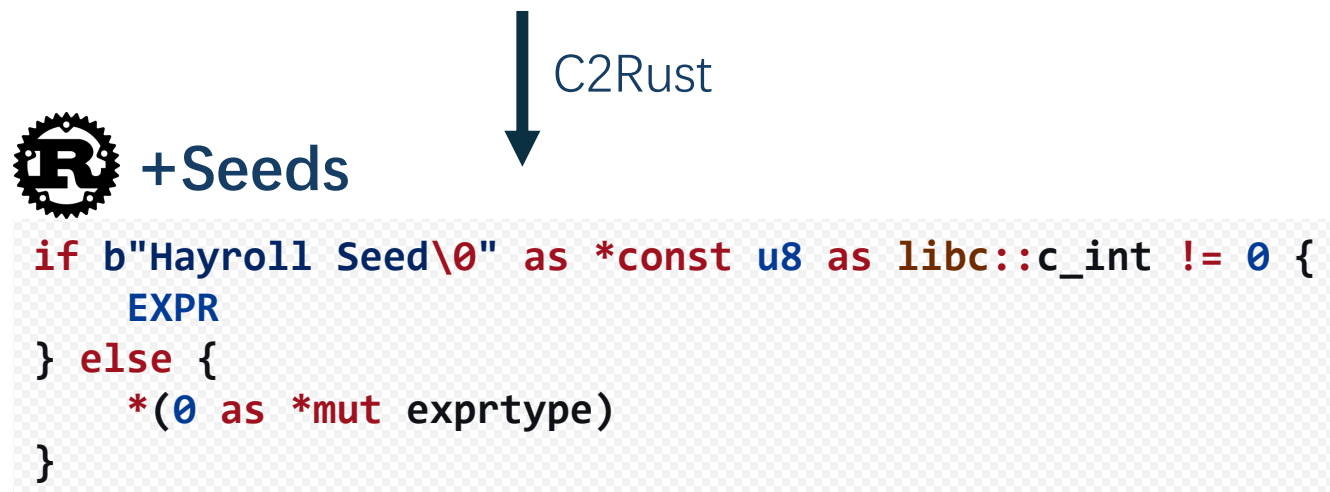


Seeding Expression



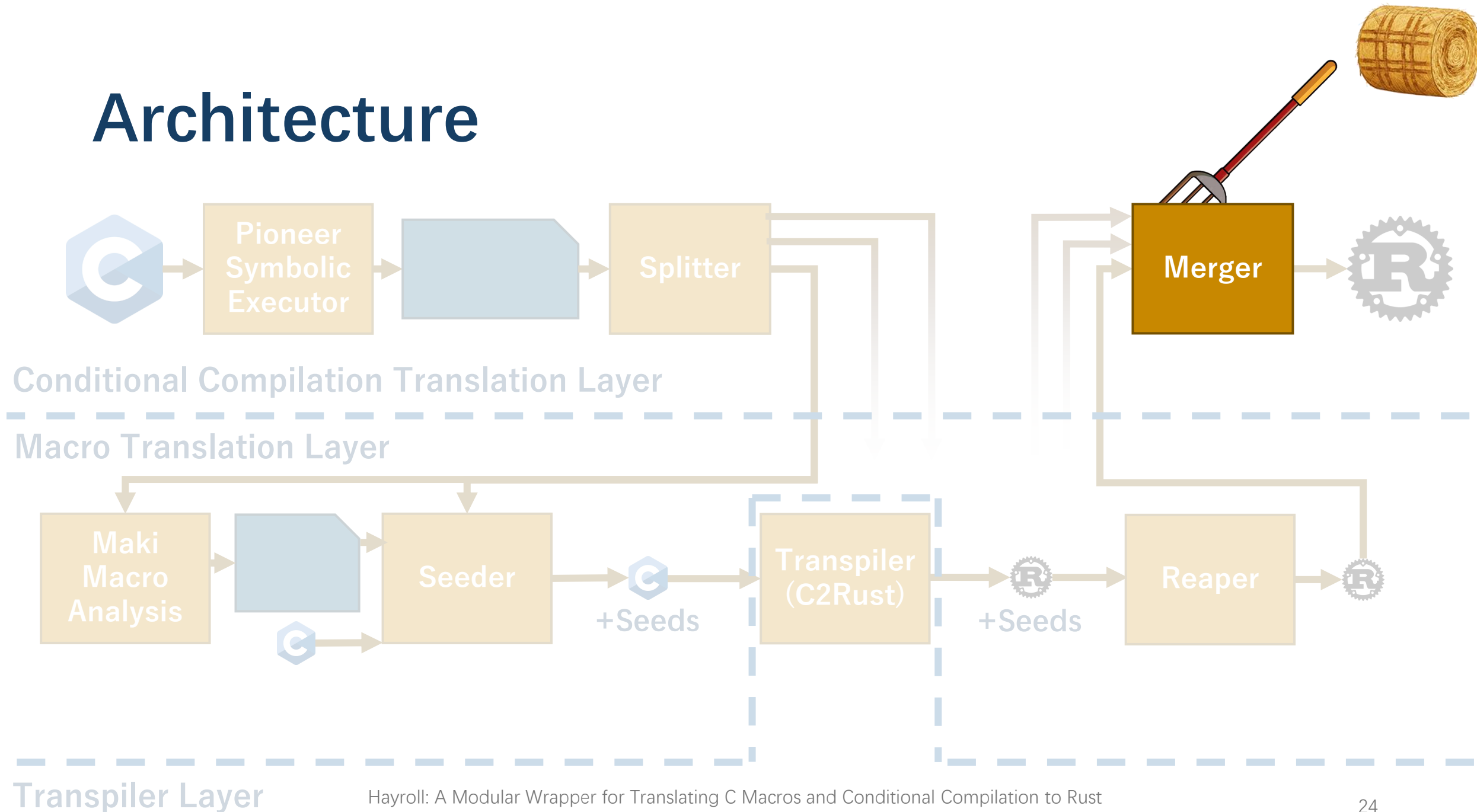
Properties of seeded AST nodes:

- Syntactically preserving
- Semantically preserving
- Carries metadata



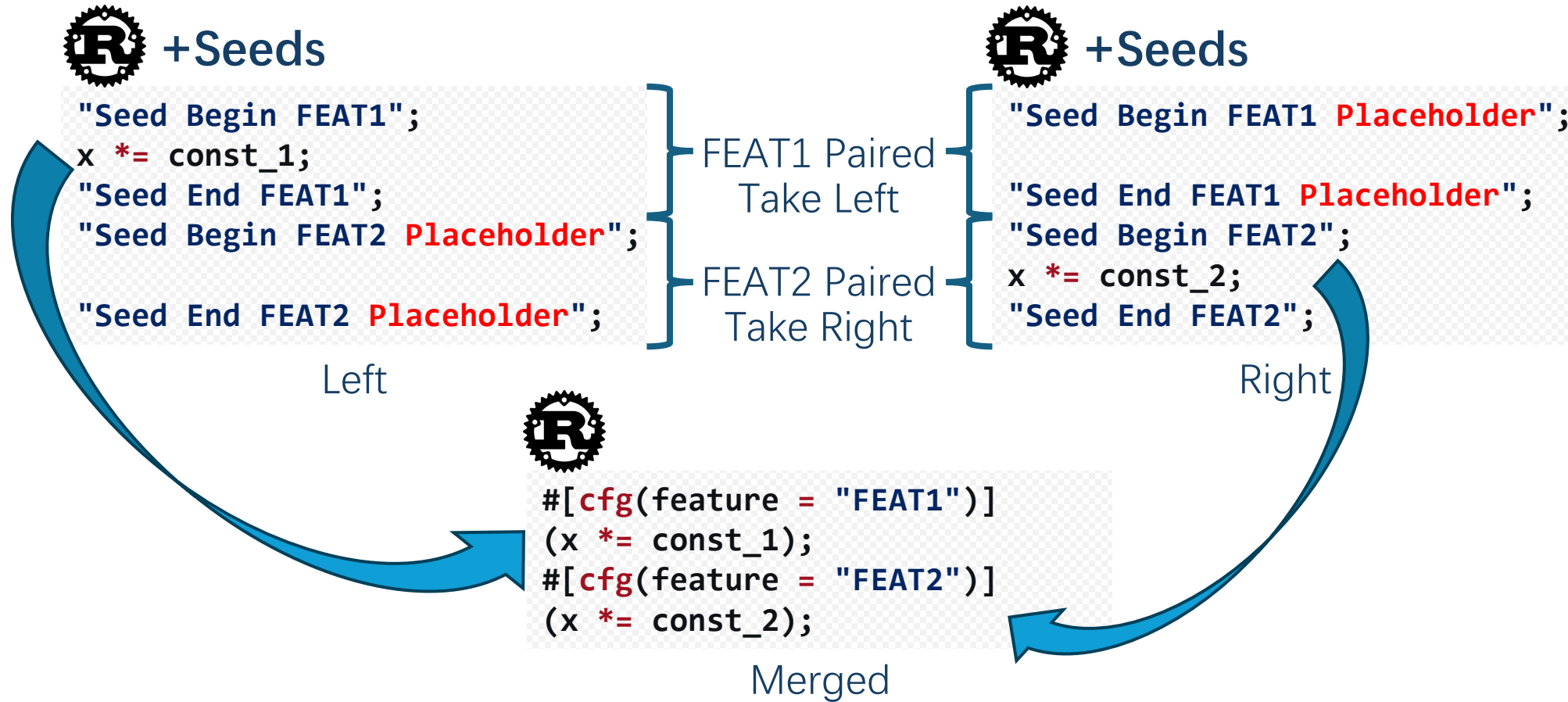
Translated into an **if expression** in Rust

Architecture





Merging with Placeholder Seeds





Hayroll Preserves Abstraction and Configurability in Transpilation

Cond Stmt Expr



```
float sinh(float x)
{
#ifdef __LIBMCS_FPU_DAZ
    x *= __volatile_onef;
#endif
    float t, w, h;
    int32_t ix, jx;
    GET_FLOAT_WORD(jx, x);
    ix = jx & 0x7fffffff;
    /* x is INF or NaN */
    if (!FLT_UWORD_IS_FINITE(ix)) {
        return x + x;
    }

    h = 0.5f;
```



```
pub unsafe extern "C" fn sinh(mut x: libc::c_float) -> libc::c_float {
    #[cfg(feature = "LIBMCS_FPU_DAZ")]
    (x *= __volatile_onef);
    let mut t: libc::c_float = 0.;
    let mut w: libc::c_float = 0.;
    let mut h: libc::c_float = 0.;
    let mut ix: int32_t = 0;
    let mut jx: int32_t = 0;
    GET_FLOAT_WORD(&mut jx, &mut x);
    ix = jx & 0x7fffffff as libc::c_int;
    if FLT_UWORD_IS_FINITE(&mut ix as *mut int32_t) == 0 {
        return x + x;
    }
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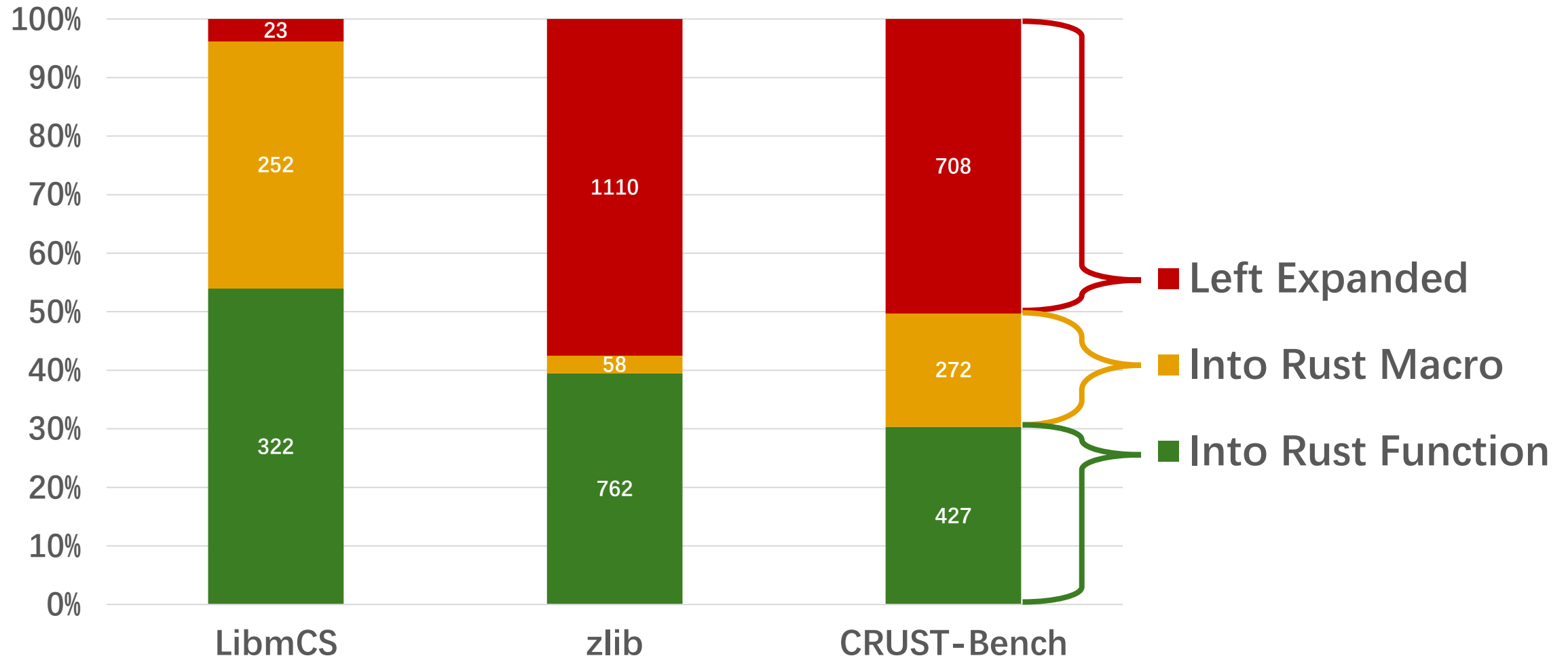


Evaluation

- **Test programs:** LibmCS, zlib, CRUST-Bench
- **Correctness:** Hayroll passes all test cases that C2Rust passes
- **Translator performance:** ~ 500 LoC/s, $\sim 20x$ slower than C2Rust
- **Coverage:** See next slide

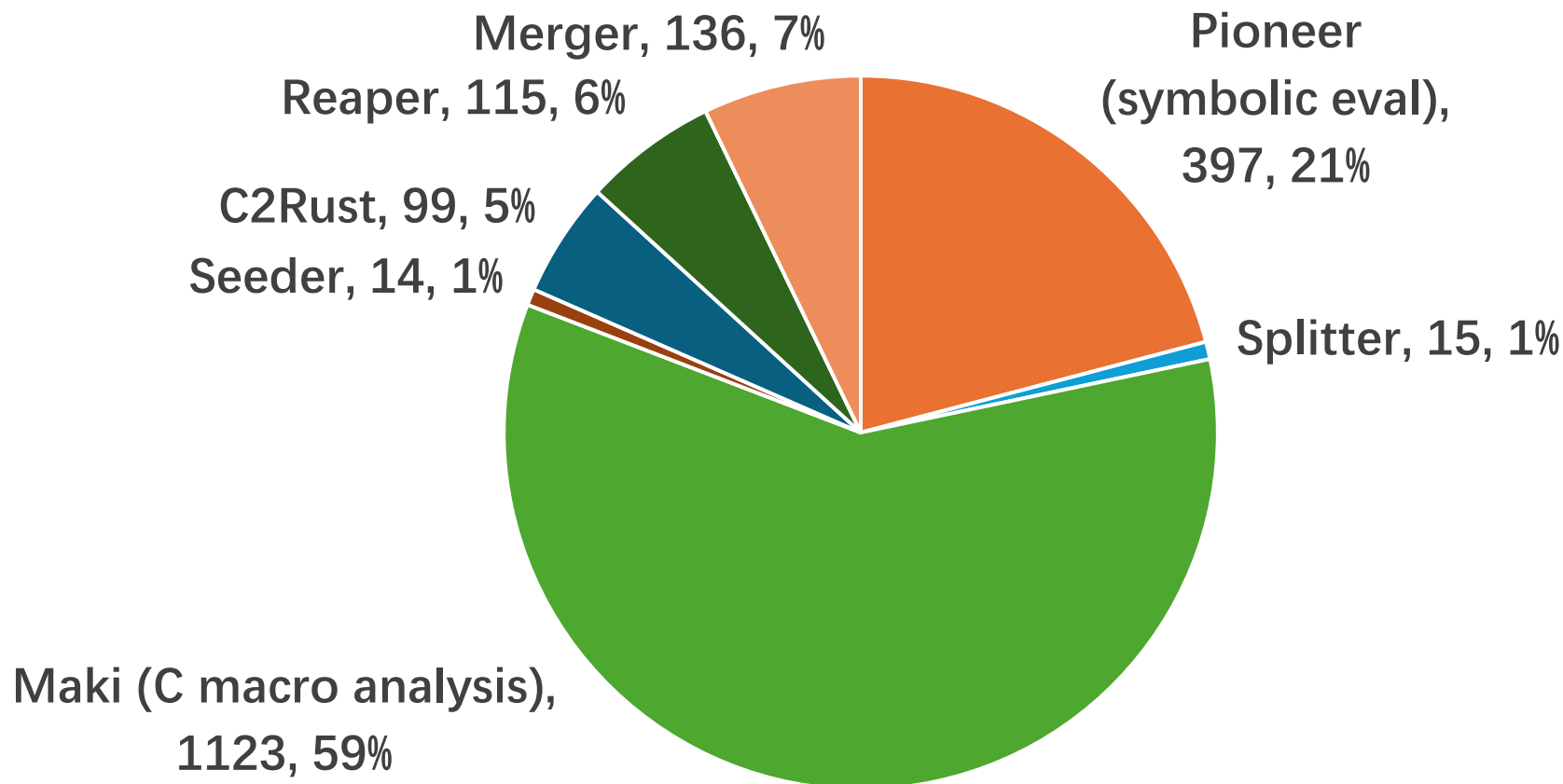


Fraction of Macros Translated



Per-component Performance (ms/kLoC)

Average over All Programs





Summary

- Hayroll translates **C macros** into **Rust functions or macros**
- Hayroll translates **C conditional compilation** into **Rust cfg attr**
- Open-sourced at <https://github.com/UW-HARVEST/Hayroll>
- The C2Rust team is working on integrating Hayroll

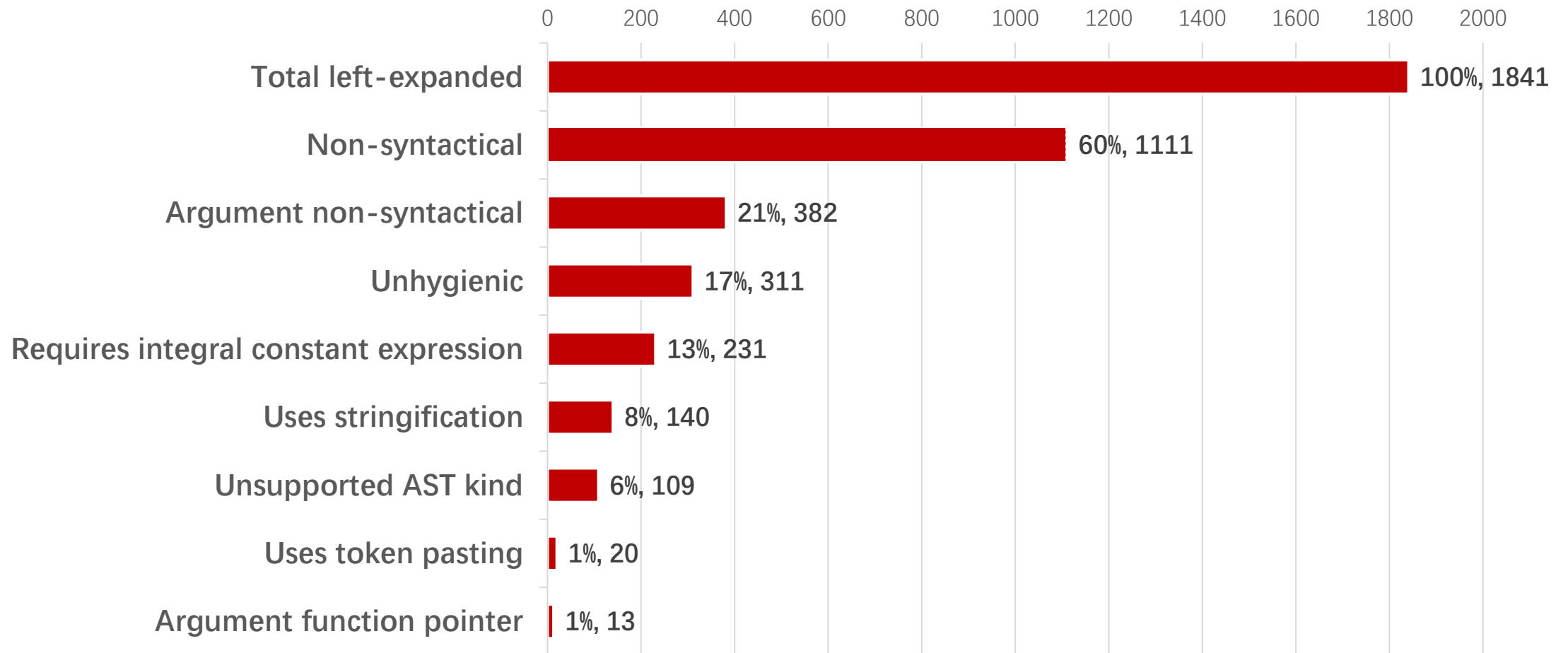


Fork us!



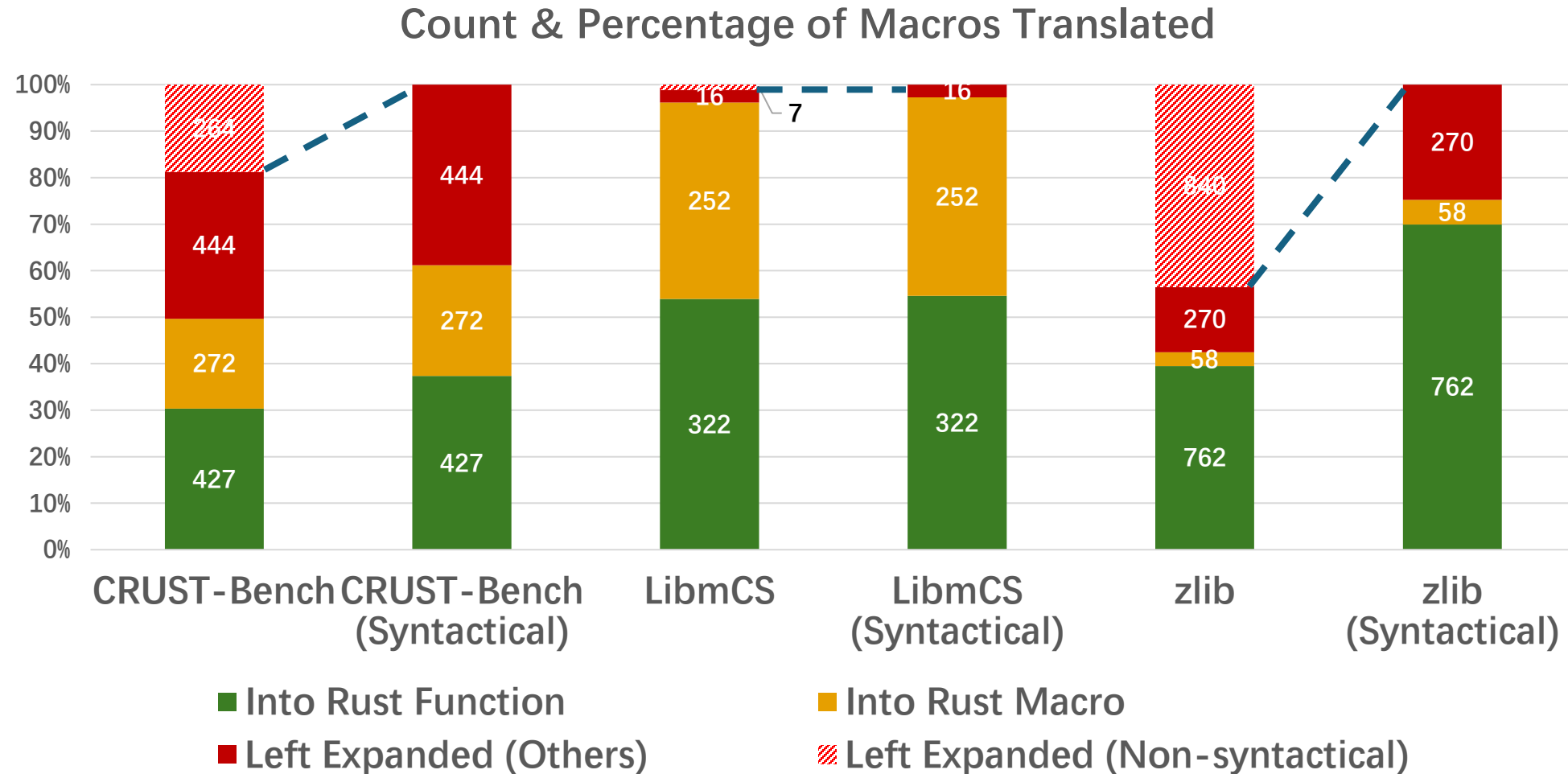


Reasons why Macros are Left Expanded





Effectiveness





Seeding Lvalue Expression



Hayroll
Seeder



```
*(  
  "Hayroll Seed" ?  
  &(LVALUE) :  
  (*(__typeof__(LVALUE)*)(0))  
)
```



C2Rust



```
*if b"Hayroll Seed\0" as *const u8 as libc::c_int != 0 {  
  RVALUE  
} else {  
  0 as *mut rvaluetype  
}
```

Lemma:

$x = y; \Leftrightarrow *(&x) = y;$



Why Translate C to Rust

- The C programming language underlies most of the world's software vulnerabilities, mostly due to **memory unsafety**
- About 70% of critical vulnerabilities in Android and Chrome are due to **memory unsafety**
- Rust is designed to be a **memory-safe** language
- A White House report mentions **Rust as a memory-safe** alternative to C/C++
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Minimap

